# Randomness 

# Randomness In Theory and Practice 

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# Randomness 



# Randomness 




## Randomness <br> Mathematics

## 1. Tournaments




0

0

Rule: every 3 players O $\begin{gathered}\text { simultaneously lose } \\ \text { to someone. }\end{gathered}$

0

0
0
$0$


0

O

Rule: every 3 players O | $\begin{array}{c}\text { simultaneously lose } \\ \text { to someone. }\end{array}$ |
| :---: |

0

0
0

$0$

0

0
0

Rule: every 3 players

O | $\begin{array}{c}\text { simultaneously lose } \\ \text { to someone. }\end{array}$ |
| :---: |

0
O
0


0

0
0

Rule: every 3 players

O | $\begin{array}{c}\text { simultaneously lose } \\ \text { to someone. }\end{array}$ |
| :---: |

0

0
0

## 0

0

## O Can this be done? O

0
0
0

## 0

0
0

## 0 <br> YES.

0

0
O
0

O

O
O

O

## Actually...

O

0
0
0

0
0
O

O No.
0

O
O
O









O Fail probability: 7/8

0


O

O
O
O

O
All

O
O

0

0

$$
\begin{aligned}
& \text { Total Failed } \\
& \text { Tournaments: } \\
& \binom{100}{3}\left(\frac{7}{8}\right)^{97}
\end{aligned}
$$

0

# At least 2/3 of those random assignments succeed. 



Method

## Sometimes it's easier to randomly choose than to explicitly construct.

# Often, 99.9\% of the objects have what you need... 

...but you can't find them!



# 2. High Girth and Chromatic Number 

## Graphs

## Graphs



Vertices

## Graphs



Edges

## Graphs



Vertices $=\{A, B, C, D, E, F\}$
Edges $=\{\{A, B\},\{A, E\},\{B, C\},\{C, D\},\{C, F\},\{B, F\}\}$

## Graphs



Cycles

## Graphs



## Cycles

## Graphs



Shortest Cycle $=$ Girth $=3$

## Graphs



Coloring

## Graphs



Chromatic Number $=3$

Can a graph have both high girth and large chromatic number?

High girth: 2-colorable locally everywhere.






# How can we force many colors if every region is 2-colorable? 

That's HARD.

Unless...


- Take many vertices
- Choose each edge with probability $p$
- Carefully choose p
- Make the graph have few short cycles and small independence number
- Remove vertices to eliminate short cycles
- Voila!

3. The Existence of Designs



7 points
Sets of size 3
Every 2 points are in exactly 1 set


# $n$ points <br> Sets of size q 

 Every $r$ points are in exactly $\lambda$ sets$n$ points
Sets of size q
Every r points are in exactly $\lambda$ sets

## nqr入?

## ngr入?

## Divisibility conditions Finitely many exceptions in $n$

Asked: 1853

Answered: Jan 15, 2014

## Solver: Peter Keevash

Method: Randomized Algebraic Construction

- Rephrase the problem as hypergraph matching
- Seek a matching by randomly picking edges and deleting their overlaps
- The beginning is easy, but the end is out of control
- Keevash: Cleverly pick stand-ins for the end game


## 4. The Crossing Lemma

# Crossing Number: 

 Minimum number of edge crossings in a plane drawing of a graph

## Crossing Lemma: if $e \geq 4 v$ then $\operatorname{cr}(G) \geq \frac{e^{3}}{64 v^{2}}$

- Fact (easy): $\operatorname{cr}(G) \geq e-3 v$
- Start with any graph
- Pick a random subgraph H by choosing each vertex with probability p
- Find the expected number of vertices, edges and crossings of H
- Apply the easy fact. Pick the best p. Done!


## 5. Algorithms

## Quicksort

## Codes

## Primality Testing

## Min-Cut

## Matrix Testing

Machine Learning


Machine Learning


## Machine Learning

Learn from data: features and outcomes
Predict: Given features, what's the outcome?

## Decision Trees



## Decision Trees are great...



## ...Random Forests are Even Better.


...Random Forests are Even Better. Train 100 decision trees with random data and random features

Merge into one predictor

## Random Forests

## Perhaps the single most successful Machine Learning paradigm

## Ever.

## Summary

## Randomness

Mathematics

## Randomness

## Mathematics

Artificial Intelligence

